

Introduction to Cellular automata

There is a wealth of literature about cellular automata, as well as many Internet resources (you'll find some of them in the [links](#) section). The aim is here much more limited. This site being devoted to laymen, I will content myself with answering both main questions any person discovering CA often ask, generally after a period of intense perplexity :

1. What might this be ?
2. What could be the applications for such a thing ?

The answer to these questions is unfortunately far from being simple. CA are abstract constructions with quite complex properties not very accessible.

A- History

The history of CA dates back to the forties with Stanislas Ulam. This mathematician was interested in the evolution of graphic constructions generated by simple rules. The base of his construction was a two-dimensional space divided into "cells", a sort of grid. Each of these cells could have two states : ON or OFF. Starting from a given pattern, the following generation was determined according to neighbourhood rules. For example, if a cell was in contact with two "ON" cells, it would switch on too ; otherwise it would switch off. Ulam, who used one of the first computers, quickly noticed that this mechanism permitted to generate complex and graceful figures and that these figures could, in some cases, self-reproduce. Extremely simple rules permitted to build very complex patterns. On that basis, the following question was asked : can these recursive mechanisms (i.e. in that case depending on their own previous state) explain the complexity of the real ? Is this complexity only apparent, the fundamental rules being themselves simple¹ ?

As a sideline, John von Neumann, relying on A. Turing's works, interested himself on the theory of self-reproductive automata and worked on the conception of a self-reproductive machine, the "kinematon". Such a machine was supposed to be able to reproduce any machine described in its programs, including a copy of itself². The most famous of his machines is the monolith of the series "2001 Space Odyssey". To change Jupiter into a star, a first monolith self-reproduces, as well as its descendants, the population so increases exponentially to quickly reach the size necessary to realize such a gigantic task.

Ulam suggested von Neumann to use what he named "cellular spaces" to build his self-reproductive machine. He could so liberate himself from real physical constraints to work in an extremely simplified universe that was nevertheless able to generate a high complexity. The use of this formal universe led him to notice : "By axiomatizing [self-reproductive automata with cellular automata], one has thrown half of the problem out the window, and it may be the more important half. One has resigned oneself not to explain how these parts are made up of real things, specifically, how these parts are made up of actual elementary particles [...] The question we hope to answer now [...] is : what are the basic principles which underlie the organization of these elementary parts in living organisms ? My discussion

will be limited to that point of view."3. On this base, he designed an about 200.000 29 states cells , containing a universal replicator, a description of itself and a Turing machine for supervision4.

CA left laboratories in 1970 with the now famous Game of Life of John Horton Conway.

B- The Game of Life

Originally, the Game of Life is presented as a mathematical game. Its description will allow us to materialize and better understand what CA are.

Like Ulam's cellular spaces, the game of life is based a grid constituted of cells, for example :

Example of a starting pattern

00	01	02	03	04
10	11	12	13	14
20	21	22	23	24

The universe is here limited to a rectangle of 5 by 3. To make the explanation easier, we numbered the cells from 0 to 4 horizontally and from 0 to 2 vertically. Light cells are active ones.

In the game of life, any adjoining cell is considered as neighbour, including diagonals.

Determination of neighbourhood

00	●	●	●	04
10	●	12	●	14
20	●	●	●	24

The graphic above shows the neighbourhood of cell 12. In this case, two cells are active out of the 8 neighbours.

The rules of the game of life are quite simple :

1. One inactive cell surrounded by three active cells becomes active ("it's born") ;
2. One active cell surrounded by 2 or 3 active cells remains active ;
3. In any other case, the cell "dies" or remains inactive.

We can interpret these rules by considering that a birth supposes a certain gathering of population, (3 in this case), that the cells cannot survive to a too wide isolation (less than two neighbours) and that a too strong concentration will kill them (more than 3 neighbours).

CA work in a *discrete* manner. That is to say time goes step by step. This means that in our case, for the g generation, each cell examines its environment and determines its future state. When all the cells have fulfilled this computation, the transitions occur. We so simulate a simultaneous treatment.

Let us illustrate this mechanism starting from the previous pattern :

First generation

1	2	3	2	1
1	1	2	1	1
1	2	3	2	1

In the previous pattern, the number of active neighbours is noted for each cell :

1. The cells 00, 04, 10, 14, 20 and 24 have got one active neighbour and then remain inactive.
2. The cells 01, 03, 21 and 23 have got two active neighbours, and then do not change.
3. The two inactive remaining cells (02 and 22) have got three active neighbours, the rule 1 is applied : they are born.
4. The cells 11 and 13 have only one active neighbour : they die.
5. Finally the cell 12 having two active neighbours, it remains alive.

For the next generation, only the cells 02,12 and 22 are then active.

Second generation

00	01	02	03	04
10	11	12	13	14
20	21	22	23	24

We show there the three fundamental properties of CA⁵ :

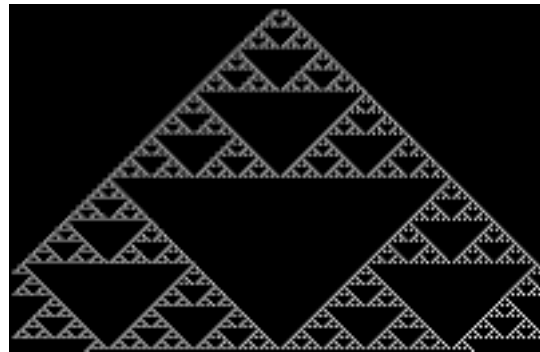
1. *Parallelism* : A system is said to be parallel when its constituents evolve simultaneously and independently. In that case cells update are performed independently of each other.
2. *Locality* : The new state of a cell only depends on its actual state and on the neighbourhood.
3. *Homogeneity* : The laws are universal, that's to say common to the whole space of CA.

C- Other Cellular Automata

The game of life is only one type of CA among an infinity. It is indeed possible to play on the whole rules that govern the universe of CA.

The most obvious parameter is the number of dimensions. Indeed, nothing obliges to consider two dimensions environments. The theoretical analysis of CA was mainly made out of one dimension automata. By reducing the number of dimensions, one limits the combinatory explosion, hence the number of possible automata. If we consider the simple case of a three cells neighbourhood, i.e. the concerned cell and its right and left neighbours, in a one dimension and two states automaton, there are only 2 power $2^3=256$ possible rules. With one dimension automata (i.e. one line) we use the second dimension to represent time. For each generation, a new line is added below the former one, we can so visualize the dynamic of this type of automata.

Example of a 1 dimension automaton (Pascal's triangle)



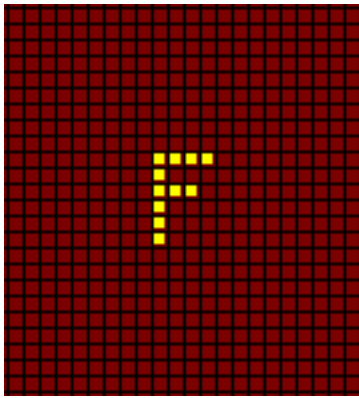
It is obviously possible to create three (or more) dimensions automata.

It is also possible to modify the determination of neighbourhood. If we consider two dimensions automata, the most common neighbourhoods are⁶ :

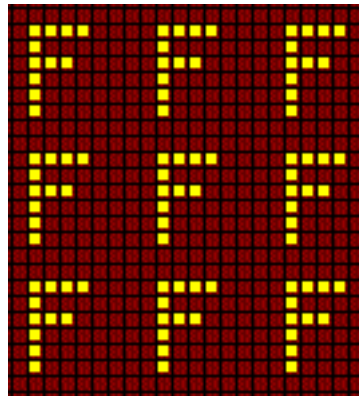
1. Von Neumann : only North/South/East/West neighbours are considered ;
2. Moore : one adds the diagonals. It's the case of the game of life ;
3. Extended Moore : one extends the distance of neighbourhood beyond one ;
4. Margolus : one considers groups of 2x2. It's this type of neighbourhood that is used in the simulation of gas behaviour.

For example, Fredkin's automata, that uses a Moore neighbourhood is based on the parity of neighbourhood. It's a *totalistic* automaton, that is to say the state of the cells depends on the sum of the states of neighbouring cells. In this case there is reproduction only if there is an odd neighbourhood value. This automata has got the remarkable property to reproduce nine copies of any basic pattern. Fredkin's rule can easily be generalized to more than two dimensions⁷.

Fredkin generation 0

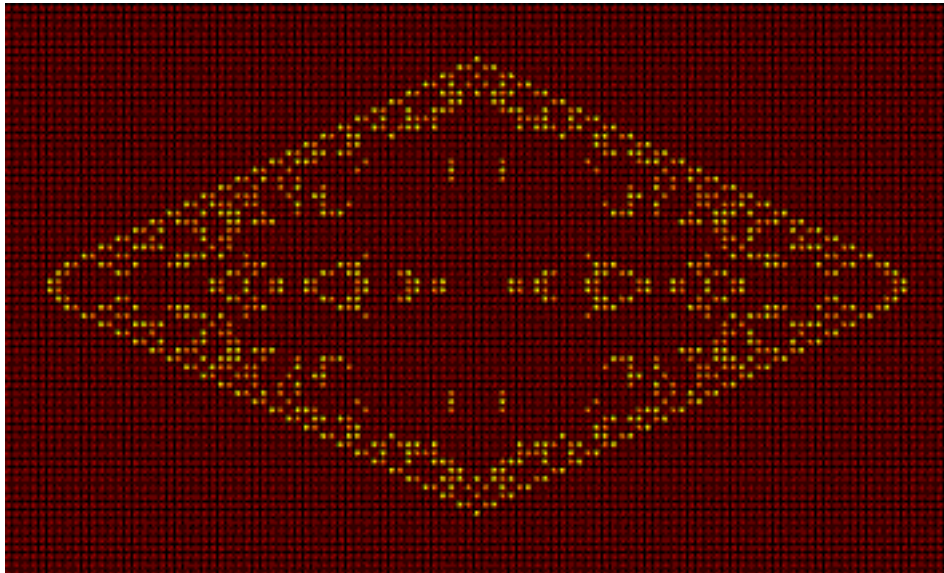


Fredkin generation 8



It is also possible to modify the number of states. You needn't restrict yourself to both states life/death. Numerous famous automata use more than two states. One of the most famous is Brian's Brains presented by Brian Silverman in 1984. This three states automaton (life, ghost, death) generates a wide diversity of complex gliders within astonishing graphic patterns.

Brian's Brain



More complex rules can be imagined. It's possible, for example, to build stochastic automata whose transition rules integrate a probability function.

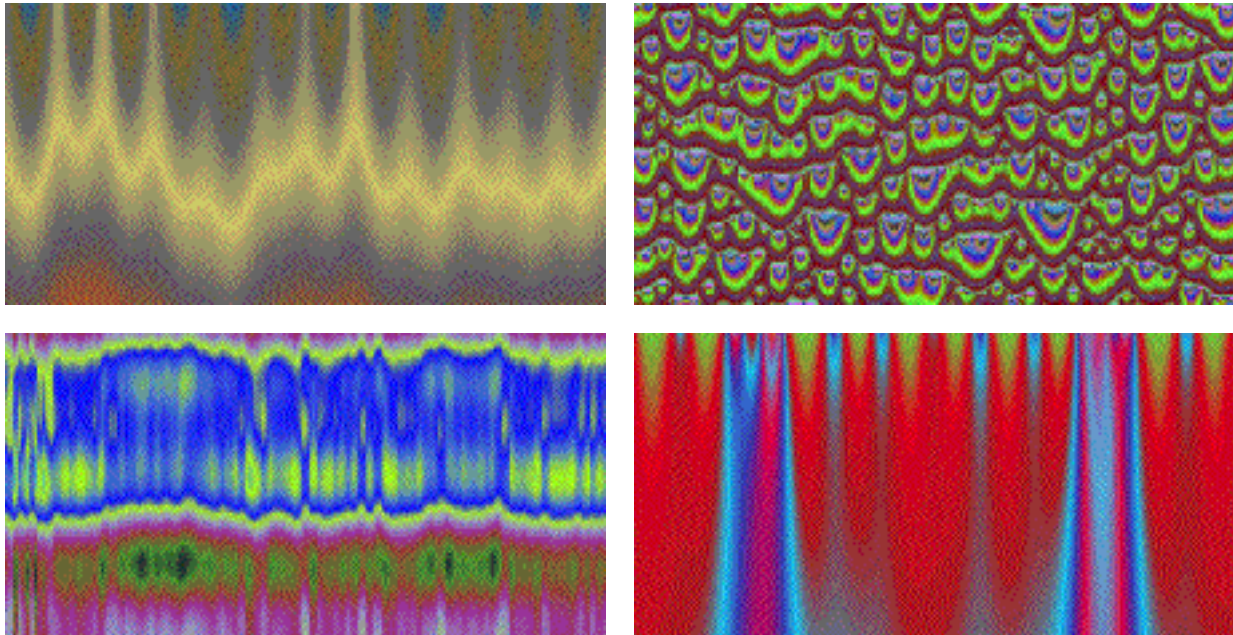
In a general way, it is possible to build any type of automata by playing on *structural* and *functional* rules. The first ones define the spatial structure of the automata network, that is its number of dimensions, the disposition of cells (squares, hexagons,... in a two dimensional automaton) and the type of neighbourhood determination. The second ones will determine the number of states and the transition rules⁸. The choice of these two types of rules permits to build a universe adapted to the demanded aim.

D- Applications

CA applications are diverse and numerous. Fundamentally, CA constitute completely known universes. Our Universe is submitted to the laws of physics. These laws are only partly known and appear to be highly complex. In a CA laws are simple and completely known. One can then test and analyse the global behaviour of a simplified universe, for example :

1. Simulation of gas behaviour. A gas is composed of a set of molecules whose behaviour depends on the one of neighbouring molecules.
2. Study of ferromagnetism according to Ising model : this model (1925) represents the material as a network in which each node is in a given magnetic state. This state - in this case one of the two orientations of the spins of certain electrons - depends on the state of the neighbouring nodes.
3. Simulation of percolation process.
4. Simulation of forest fire propagation.
5. In a different field, CA can be used as an alternative to differential equations⁹.
6. Conception of massive parallel computers.
7. Simulation and study of urban development¹⁰.
8. Simulation of crystallisation process.
9. (...)

In a more daily field, CA can be used as graphic generators¹¹. The several images below, generated with [Capow](#), show some graphic effects.



But it is undoubtedly in the field of artificial life that CA are most known.

E- Cellular Automata and Artificial Life

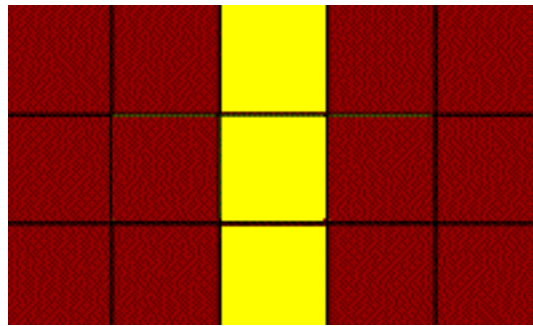
Since the beginning, von Neumann, considered CA as tools to analyse the reproduction mechanisms of

living being. Indeed, the properties of CA permit to show and analyse some of the Living fundamental mechanisms.

1- Emergence

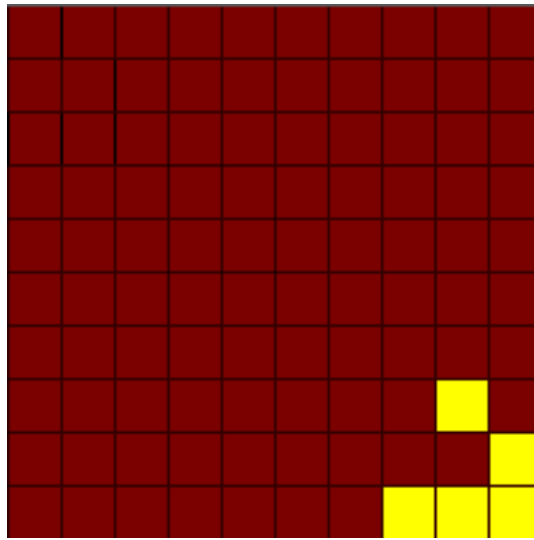
The notion of emergence appeared with the general systems theory¹². In a general way it is said that "*the whole is greater than the sum of the parts*" or "*The global behaviour is greater than sum of the behaviours of the individual parts*". These expressions mean that the complex association of elements induces the apparition of new phenomena and mechanisms. "At each level (of the prebiotic, biotic and social evolution) new properties appear that cannot be explained by the sum of the own properties of each part that constitute the whole. There is a qualitative gap (...) the property of emergence is linked to complexity. The increase of the diversity of elements, the increase of the number of links between these elements, and the game of non linear interactions lead to hardly predictable behaviours."¹³. The so-called "global" emergence then characterize the properties of a system that are new ones comparing to the properties of its isolate components, or organized in a different way. Life is undoubtedly part of these¹⁴.

We examined previously the behaviour of a line of three vertical cells : at the first generation, we get three horizontal cells, at the second one, three vertical cells again. A line of three living cells then generates a cycle.



This pattern is a "blinker". A blinker is not made up of a group of given cells, it is a *dynamic* pattern that emerges within the space of CA. The blinker seems to be independent ; it is a specific structure, particular inside its environment.

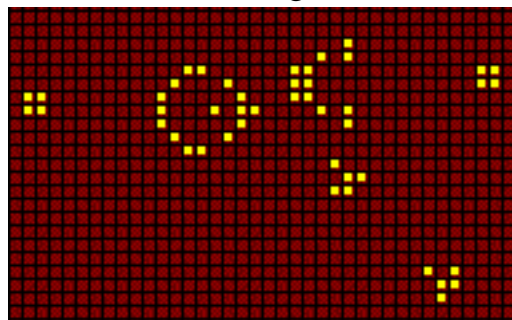
The rules of the game of life were determined so as to generate a wide diversity of unpredictable structures. The specialists listed a whole faun of patterns whose behaviours are astonishing. Complete libraries are available (see [links](#)). One of the most famous is the glider. A given five cells pattern self-reproduces every four generations one cell away.



Even more than blinkers, gliders evoke emergence. It looks like a creeping animal, going all over the space on straight line. A glider is not a set of cells. At each generation, its cells are replaced. As the atoms that constitute you are not the ones you had at your birth, the components of the glider are constantly renewed. The application of the rules of the game of life induce the apparition of a dynamic, coherent and independent structure, with peculiar properties : this is the very nature of emergence. These properties can be used for specific aims. The glider is used to represent a signal, you will find an example in [LogiCell](#).

Another remarkable pattern is also to be noted : the *glider gun*. It is a set of cells generating gliders. It made it possible to demonstrate that the population of the game of life could grow indefinitely. The classical p30 gun is used as a generator in [LogiCell](#).

Glider gun



2- Self-Reproduction

How far can these emergences go ? Is the apparition of life possible inside the universes of Cellular Automata ?

The idea that it should be possible to create life inside a computer is based on the computing theory proposed by Turing. The ability of a universal Turing machine to emulate any other machine - i.e. fundamentally the fact that all sort of computers are equivalent - led von Neumann to consider how automata could self-reproduce¹⁵. He more particularly wondered "what kind of logical organization of an automaton is sufficient to produce self-reproduction"¹⁶. It is the *kinematon* we mentioned above.

Some CA can be used to construct universal Turing machines. The game of life is part of them. Considering the physical impossibility to realize the kinematon, we saw that von Neumann turned to the CA to build his self reproductive automaton. This one was extremely complex since von Neumann considered a universal self replicator. In 1968, Edgar Codd proposed a simplified version of von Neumann's automata, that only used height states, but in this case again, Codd was considering construction universality. Things changed in the eighties with Christopher Langton. Langton considered that to study living systems inside a computer one only needs to consider *necessary* elements not *sufficient* one¹⁷. He then gave up the idea of a universal replicator.

Langton's basic idea is that it is possible to conceive a CA supporting a structure whose components constitute the information necessary to its own reproduction. This structure is then both itself and representation of itself.

Langton's automaton¹⁸ uses height states and twenty-nine rules. The structure that reproduces itself is a loop constituted of a sheath within which circulates the information necessary to construct a new loop i.e. necessary to reproduction.

A Langton's loop

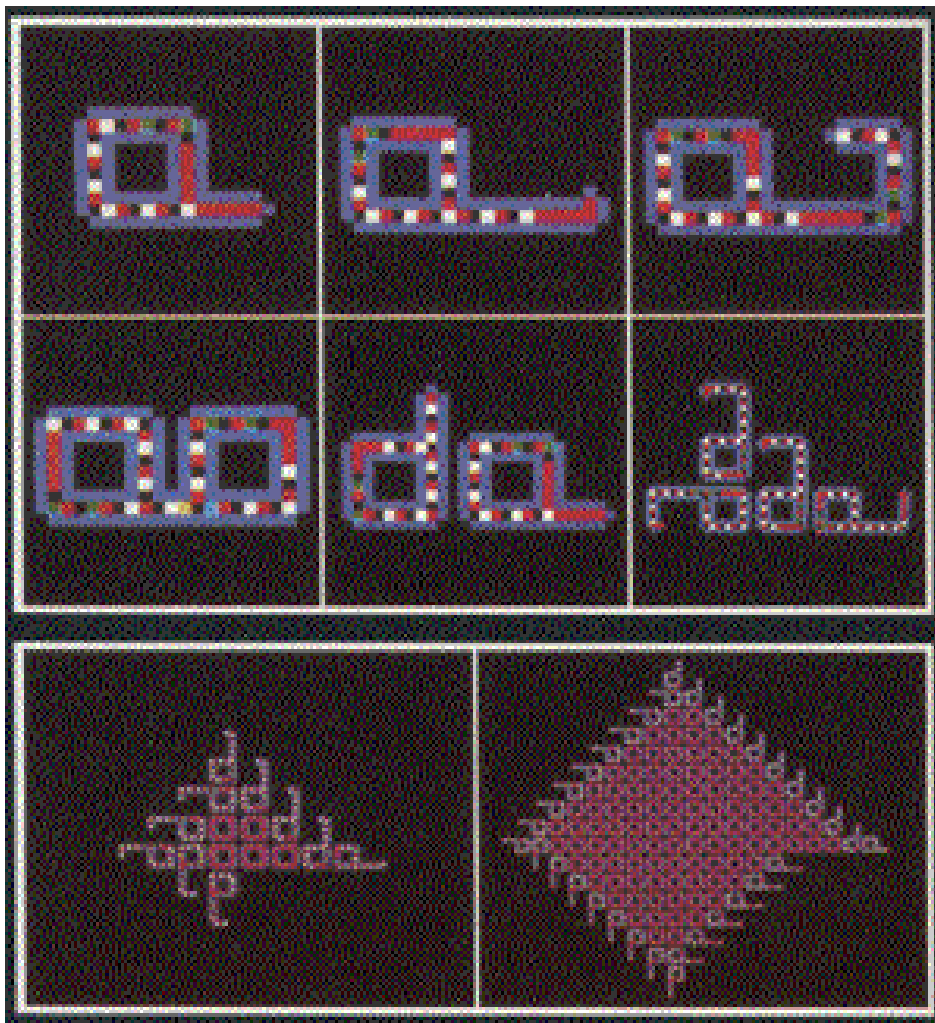
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      2 2 2 2 2 2 2 2
2 1 7 0 1 4 0 1 4 2
2 0 2 2 2 2 2 2 0 2
2 7 2           2 1 2
2 1 2           2 1 2
2 0 2           2 1 2
2 7 2           2 1 2
2 1 2 2 2 2 2 2 1 2 2 2 2
2 0 7 1 0 7 1 0 7 1 1 1 1
      2 2 2 2 2 2 2 2 2 2 2 2

```

The cells in state 2 constitute the sheath, inner cells contain the information for reproduction. They are, in some way, the DNA of the loop. The sequences 7-0 and 4-0 propagate toward the tail. When they reach the extremity, the first ones extend the tail, the second ones construct a left-hand corner¹⁹. The addition of a "sterilization" rule that blocks the evolution after a certain number of generations leads to the construction of some kind of coral.

Langton's loops



S. Levy, *Artificial Life.*, Penguin, 1992.

Like von Neumann's automata, Langton's loops show that "one of the fundamental properties of living organisms, self reproduction, can be explained in terms of interactions of simple elements and that it can be studied in its logical principles independently of its physical realization"²⁰.

Langton's loops cannot, in any way, be considered as "living", they are nothing but a limited self reproductive construction. According to Conway, if a large enough cellular space could be considered, life forms might appear within it, but it has been reckoned that it might only happen in a space of at least 10^{12} cells²¹. As a reminder, the number of particles in the universe is estimated to less than 10^{100} .

If apparition of life is envisaged, if some automata have computation universality properties, it is because these automata belong to a special category : *CA at the edge of chaos*.

3- Chaos and complexity

The number of possible CA is potentially infinite. In this context, Wolfram wondered about the existence of CA's general behaviour rules²².

S. Wolfram studied one dimensional automata, with two states and a neighbourhood of 2. He only considered as "legal" the automata that firstly eliminate any cell without neighbour, and secondly, are symmetric. There are then only 32 "legal" automata that Wolfram systematically studied.

This study showed that, according to the author, numerous CA, and maybe all of them, fall into four basic classes.

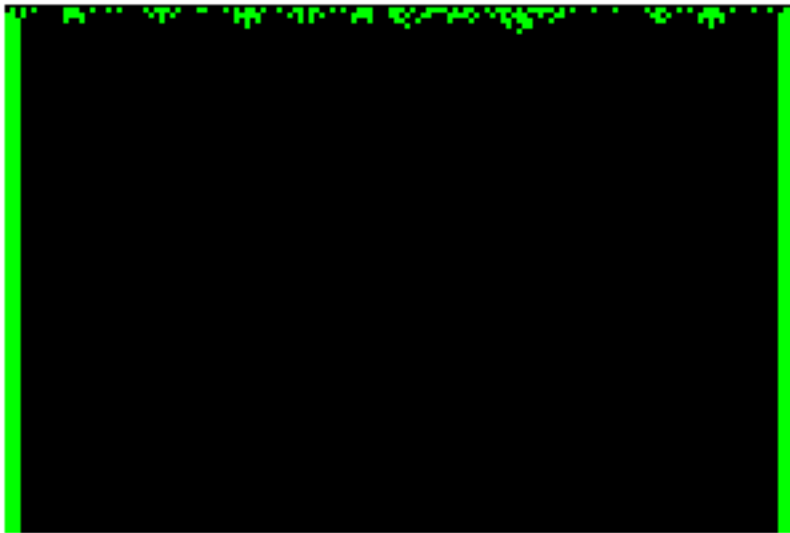
Class I- evolution leads to an homogeneous state.

Class I automata (rule 36)



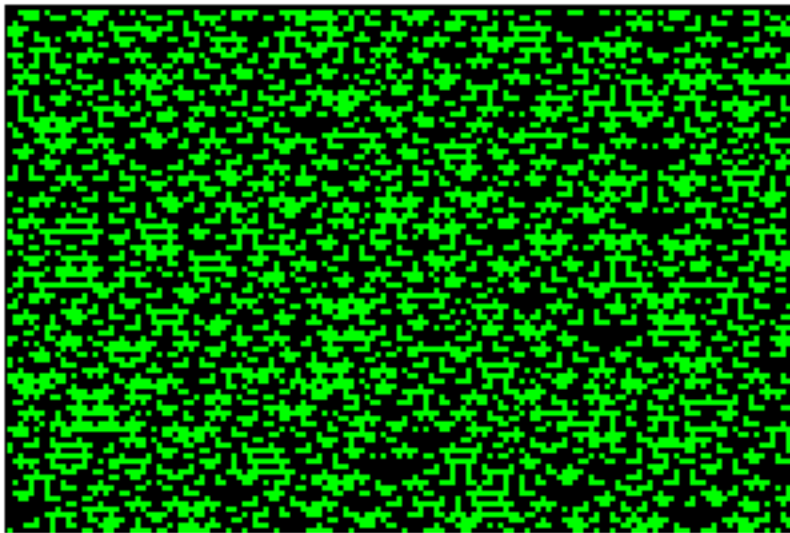
Class II- Evolution leads to simple or periodic structures.

Class II automata (rule 40)



Class III- Evolution leads to chaotic states.

Class III automata (rule 18)



Class IV- Evolution leads to complex global structures.

Class IV automata (rule 20)



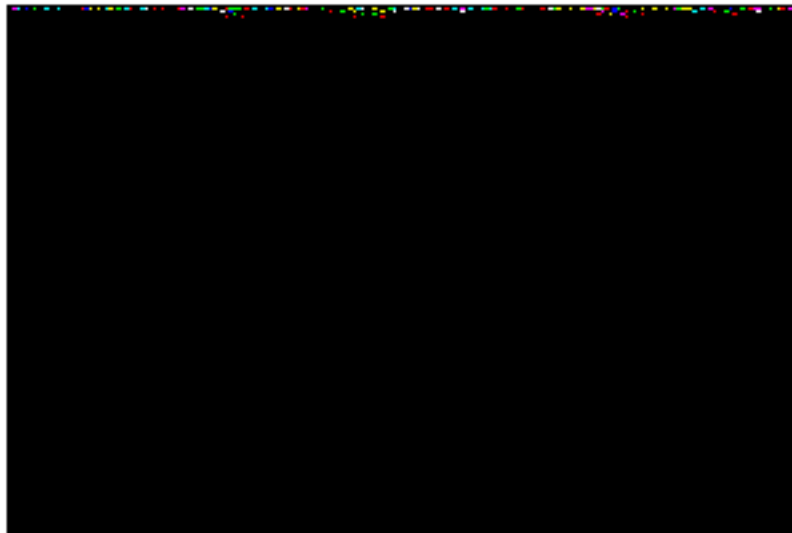
According to Wolfram, CA belonging to class IV generate structures that are strongly reminiscent of the game of life. Yet it is known that this CA has the computational universality property. This means that : "Cellular Automata may be viewed as computers, in which data represented by initial configurations is processed by time evolution. Computational universality implies that suitable initial configurations can specify arbitrary algorithmic procedures"²³. He then puts forward the hypothesis that class IV characterizes the automata having universal computation capability. In order to let this capability emerge, the cells must be able to communicate and to transmit information. In classes I and II automata cells are too strongly interdependent to allow a useful information processing. Class III automata are characterized by a too weak cells interdependence. Most of automata belong to classes I to III. They represent 30 over the 32 Wolfram's automata. CA belonging to class IV are then only to be found in a minority of cases. The CA that are at the limit between classes I and II on the one hand, and class III on the other hand, are the only ones to be capable to deal with information in a useful way, and therefore are the only "interesting" ones²⁴.

C. Langton was also interested in the existence of general classification rules of CA but in a more general way. The difficulty lies into the number of possible automata. If we consider the only automata

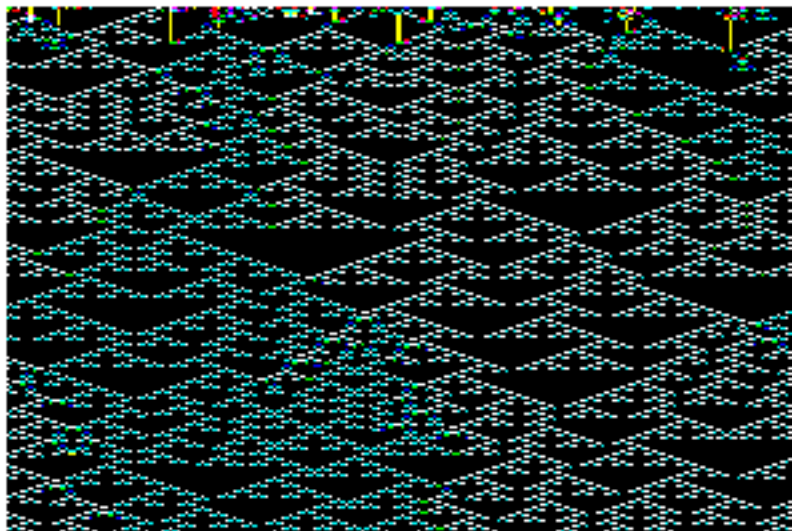
with 8 states and a neighbourhood of 5, there are 8^5 , i.e. 8^{32768} possible universes. Langton decided to classify the CA depending on a general parameter λ . The λ parameter is, in fact, the probability, within all possible neighbourhood configurations, that one given configuration should lead to an "active" cell, i.e. : $1 - (\text{number of quiescent transitions} / \text{total number of transitions})$. Using this parameter to build transition rules, Langton managed to classify CA. For a low value, cells quickly disappear. If the value is raised over about 0.2, cyclical or persisting structures appear. Over about 0.3, complex and unpredictable behaviours appear. Finally, over about 0.5, the multiplication of structures induces a chaotic behaviour. To a certain extent the λ parameter indicates the temperature of the universe of the CA²⁵.

The four following figures show examples of one dimension automata with eight states and a neighbourhood of 5²⁶.

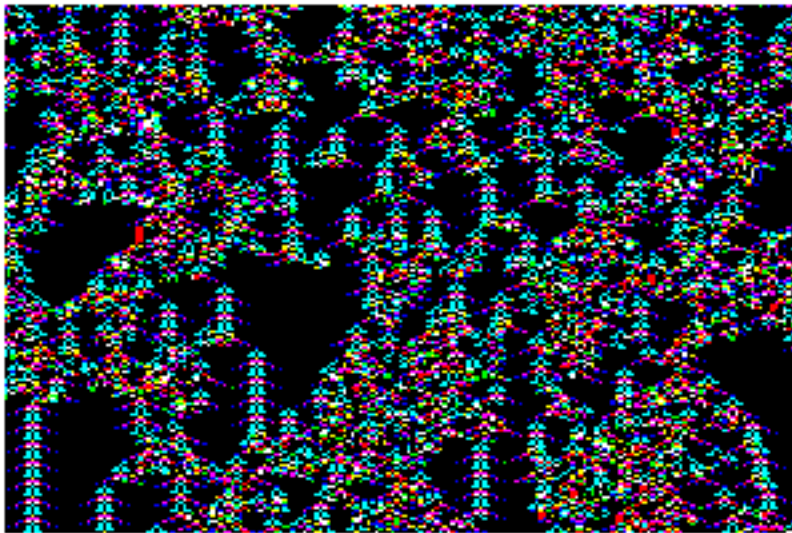
$\lambda=0.1$



$\lambda=0.25$



$\lambda=0.45$



$\lambda = 0.8$



We then find again Wolfram's classification, but in a more general framework. According to Langton, the automata belonging to class IV, those whose λ parameter is roughly to be found between 0.3 and 0.5, are those whose ability to process information is the most important. "(...) CA capable of performing non trivial computation - including universal computation - are most likely to be found in the vicinity of "phase transitions" between order and chaos (...) "[27](#).

J.C. Heudin^{[28](#)} uses a parameter (β) that, in two dimensions CA, corresponds to the number of neighbours necessary for a cell to remain unchanged. In the game of life, this parameter value is 2. He then redefines automata with four rules :

1. R1 : the neighbourhood is inferior to β : the cell dies.
2. R2 : a cell surrounded by β living cells keeps its state.
3. R3 : a cell that has $\beta + 1$ living neighbours becomes alive.
4. R4 : a cell surrounded by more than $\beta + 1$ neighbours dies.

For $\beta=1$, the probabilities of rules execution are $R4 > R3 > R2 > R1$. For $\beta=2$, we've got a complete

inversion with $R1 > R2 > R3 > R4$. Beyond 2, the order remains identical, but the rules R2, R3 and R4 are very seldom executed. Heudin then shows a fundamental change in the universe of CA for $\beta = 2$. It is during this deep modification of the automata properties - this *phase transition* - that complexity appears : "To appear [complex] needs order and a pinch of chaos. This situation occurs only at the interface of both systems, at the border who leads to chaos."²⁹.

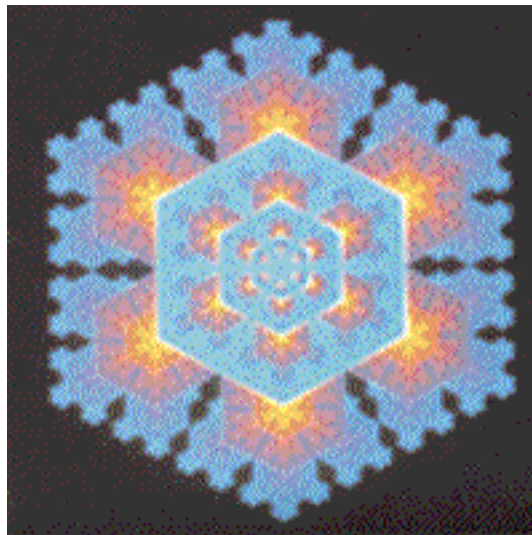
With the generalization of CA behaviours, diversity of universal structures or the existence of life, tends to show that the laws of our Universe are precisely at the border between order and chaos. This is what Langton means in the expression : "life at the edge of chaos"³⁰.

Conclusion

CA are abstract structures which make it possible to study virtual completely known universes. They help us to understand our Universe : "Thus complex physical and biological systems may lie in the same universality classes as the idealized mathematical models provided by cellular automata. Knowledge of cellular automaton behaviour may then yield rather general results on the behaviour of complex natural systems."³¹.

The universal computation capability of some CA, i.e. their capability to construct arbitrary complex structures, is not sufficient to prove the possibility of life apparition in these universes. Generalization of CA behaviours to our Universe is not direct and its pertinence is hard to be proved. However, and as a conclusion, I let you think about the two following images.

Norman Packard's snowflake



S. Levy, *Artificial Life*, Penguin, 1992.

A real shell texture and its equivalent generated with a CA



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Wolfram S., Universality and complexity in cellular automata, *Physica D*, 10:1-35, 1984. Available at : <http://www.stephenwolfram.com/publications/articles/ca/84-universality/index.html>

1- Heudin JC, *La Vie Artificielle*, Hermès, Paris, 1994, p. 35.

2- This causes a big recursion problem. The machine has to contain a self-description, but in order to be complete this description must be described too... etc... To solve this problem, the machine must be able to interpret the description as a program, a set of instructions, and as a component. The description will then be computed to construct the new machine, and then only copied in order to give the new machine a self-description. This mechanism corresponds to the current interpretation of the functioning of DNA discovered after von Neumann's work.

3- Von Neumann J. et Burks A. ed., *Theory of Self-Reproduction Automata*, University of Illinois Press, 1966, p. 77, in Ostolaza J.L., Bergareche A.M., *La vie artificielle*, Seuil, Paris, 1997, pp. 37-38. Translated from French.

4- Heudin, *Idem*, p. 39.

5- Rucker R., Walker J., *Introduction to CellLab*, <http://www.fourmilab.ch/cellab/>

6- Shatten A., *Cellular Automata*, Institute of General Chemistry Vienna University of Technology, Austria, 1997.

7- Rucker R., Walker J., *op.cit.*

8- Langlois A., Phipps M., *Automates cellulaires. Application à la simulation urbaine*. Hermès, Paris, 1997. p. 17.

9- Toffoli T., *Cellular automata as an alternative to Differential equations*, in *Modelling Physics*, *Physica* 10D, 1984.

10- Langlois A., Phipps M., *op.cit.*

11- According to Rucker et Walker *op. cit.*, "in five years you won't be able to watch television for an hour without seeing some kind of CA".

12- Bertalanffy (von) L., *Théorie générale des systèmes*. Dunod, Paris, 1973.

13- Rosnay (de) J., *Le microscope. Vers une vision globale*. Points, Seuil, Paris, 1975, pp.131-132. Translated by me.

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- 16- Shatten *op. cit.*
- 17- Adami Ch., *idem*, p. 27.
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